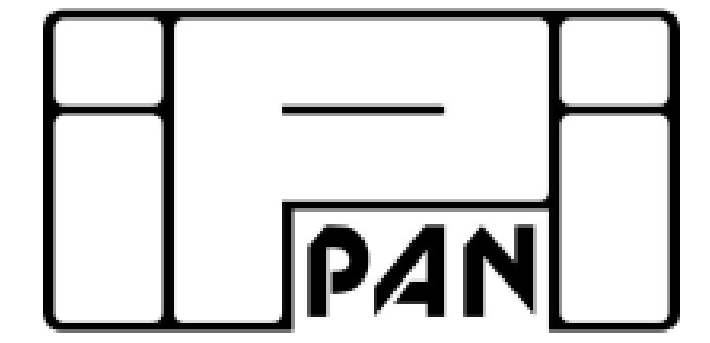


# Towards Partial Order Reductions for Strategic Ability

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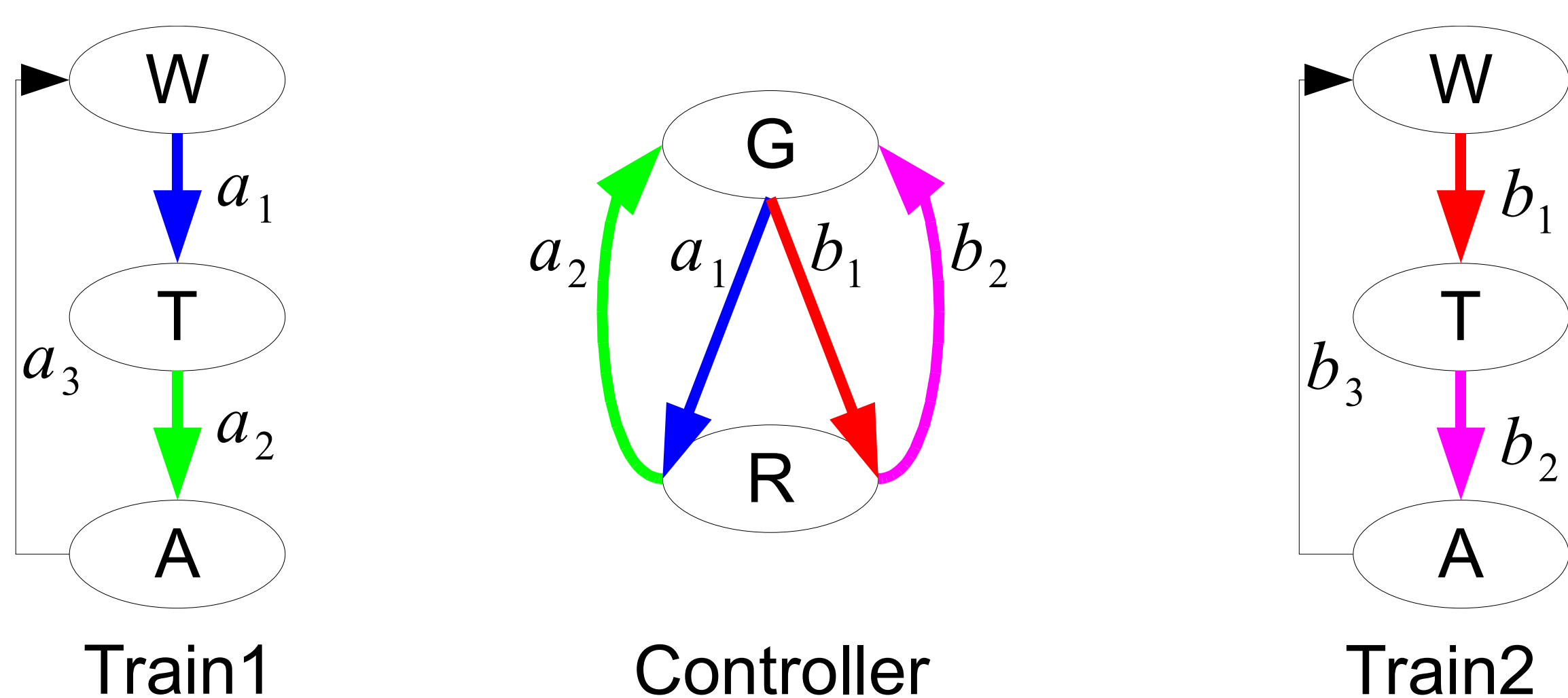


## Contribution in a Nutshell 1

- ▶ We propose a general semantics for **strategic abilities** in **asynchronous systems**, with and without perfect information
- ▶ We establish the complexity of model checking  $\sim$  no surprises here
- ▶ Most importantly, we develop a methodology for **partial order reduction (POR)** in verification of abilities under imperfect information
- ▶ The methodology is **as powerful as the reductions for linear time logic**
- ▶ Interestingly, it does not work for abilities under perfect information

## Strategic Logics 2

- ▶ **Strategic logics** provide powerful tools to reason about multi-agent systems
- ▶ We focus on a variant of **ATL\*** with imperfect information and memoryless strategies, called **sATL\*<sub>ir</sub>**
- ▶ Important for specification and verification of functionality and security requirements
- ▶ Example formulae:
  - ▶  $\langle\langle \text{alice} \rangle\rangle F \text{ win}$ : Alice can eventually win no matter what the other agents do
  - ▶  $\langle\langle \text{alice, bob} \rangle\rangle G \text{ safe}$ : Alice and Bob together can guarantee that the system will always remain in a safe state
  - ▶  $\neg\langle\langle \text{tr}_1, \text{tr}_2 \rangle\rangle F(\text{in}_1 \vee \text{in}_2)$ : trains 1 and 2 cannot enter the tunnel on their own, even if they collaborate



Asynchronous MAS: 2 trains and a controller

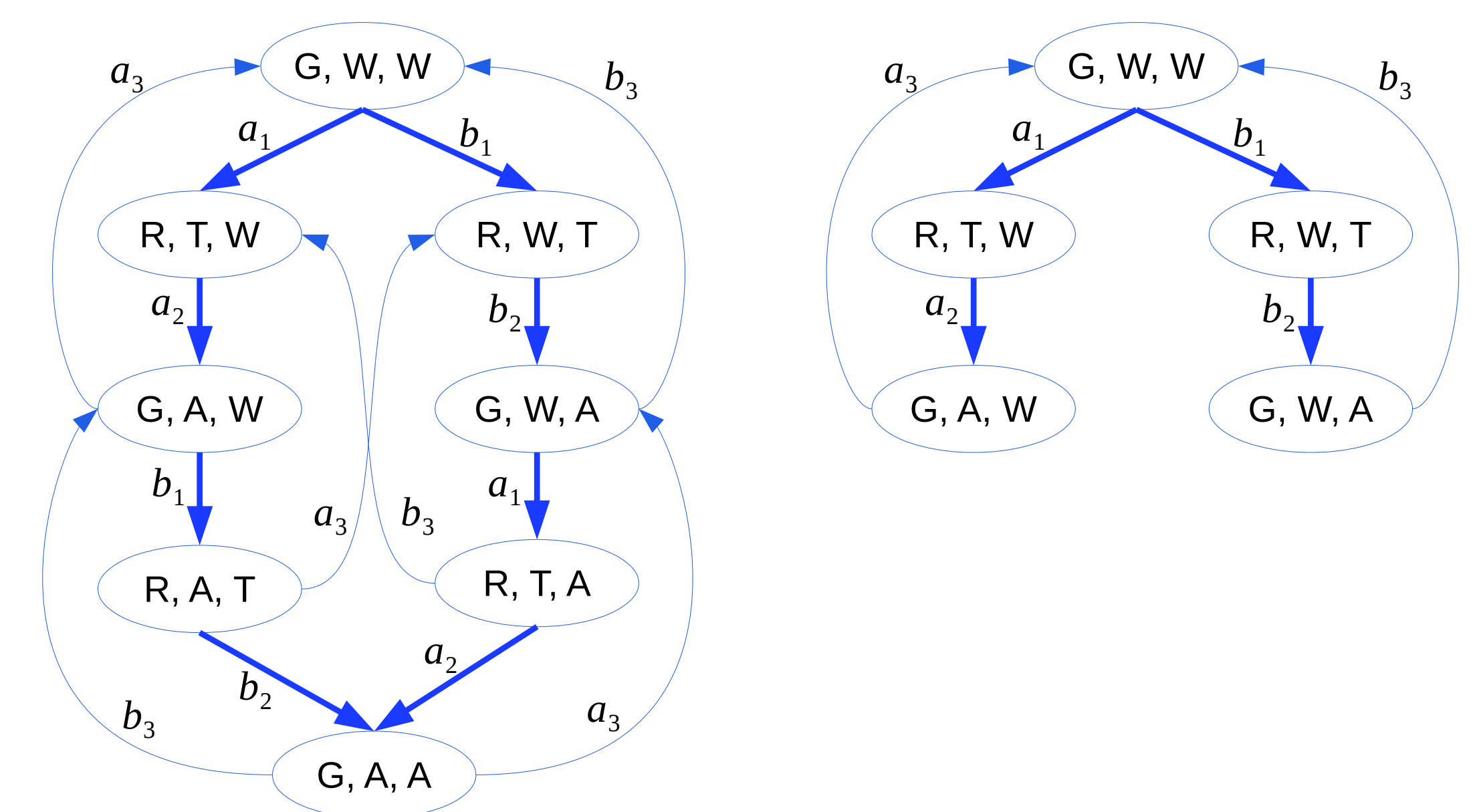
## Asynchronous Multi-Agent Systems 3

- ▶ The semantics of strategic logics are almost exclusively based on synchronous concurrent game models
- ▶ However, many multi-agent systems are inherently **asynchronous** or have simple asynchronous abstractions  $\sim$  need to adapt the semantics of strategic ability
- ▶ Moreover, asynchronous models suffer from **state-state space explosion** due to interleaving...
- ▶ ...but **effective reduction methods** exist for linear time temporal properties in asynchronous distributed systems

Question: Can we adapt the methods to the more expressive language?

## Partial Order Reductions for LTL and sATL\*<sub>ir</sub> 4

- ▶ A stack represents the visited path  $\pi = g_0 a_0 g_1 a_1 \dots g_n$  of  $M'$ .
- ▶ For  $g_n$ , the following three operations are computed in a loop:
  - ▶ The set  $en(g_n) \subseteq Act$  of enabled actions is identified and a subset  $E(g_n) \subseteq en(g_n)$  of **necessary actions** is heuristically selected.
  - ▶ For any action  $a \in E(g_n)$  compute the successor state  $g'$  of  $g_n$  such that  $g_n \xrightarrow{a} g'$ , and add  $g'$  to the stack.
  - ▶ Recursively proceed to explore the submodel originating at  $g'$
- ▶ Remove  $g_n$  from the stack.



Unfolding the trains and controller system: full model  $M$  (left), reduced model  $M'$  (right). Visible transitions are depicted by blue bold arrows

## Conditions for selection of $E(g)$ for LTL 5

- C1** Along each path  $\pi$  in  $M$  that starts at  $g$ , each action  $a \in Act \setminus E(g)$  that is **dependent** on an action in  $E(g)$  cannot be executed in  $\pi$  without an action in  $E(g)$  is executed first.
- C2** If  $E(g) \neq en(g)$ , then each action in  $E(g)$  is **invisible**, i.e., does not change valuation of  $g$ .
- C3** For every cycle in  $M'$  there is at least **one node**  $g$  in that cycle for which  $E(g) = en(g)$ .

## Theorem (1)

Partial order reductions under **C1** and **C3** preserve **sATL\*<sub>ir</sub>** under concurrency-fairness.

Assume that the actions of the agents in  $A \subseteq Agents$  are visible.

## Theorem (2)

Partial order reductions under **C1**, **C2**, **C3** preserve the formulas of **sATL\*<sub>ir</sub>** that refer only to coalitions  $A' \subseteq A$ .

## Takeaway Message 6

Takeaway message: **free lunch is sometimes possible** 😊

Main contribution: **proving** that the powerful variant of POR for LTL can be applied to **much more expressive language of sATL\*<sub>ir</sub>** (nontrivial!)